

# Substitution on infinite alphabets and generalized Bratteli-Vershik models

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# Plan for this talk

Cantor d.s

Borel d.s

- Introduction to Bratteli diagrams and generalized Bratteli diagrams
- Two versions of Kakutani-Rokhlin tower construction for a class of substitution dynamical systems on a **countably infinite** alphabets (known as left determined substitution).
- Bratteli-Vershik (B-V) models for such substitution dynamical systems.
- Using the Bratteli-Vershik model we find explicit expressions for invariant and ergodic measures for such substitution dynamical system.
- Example. ✓

# Preliminaries

- Let  $\mathcal{A}$  be a countably infinite set, called an alphabet. We denote by  $\mathcal{A}^{\mathbb{Z}}$  the bi-infinite sequence  $(x_i)_{i \in \mathbb{Z}}$  on  $\mathcal{A}$ .
- Note that  $\mathcal{A}^{\mathbb{Z}}$  is a non-compact Polish space.
- For  $x \in \mathcal{A}^{\mathbb{Z}}$  we denote by  $\mathcal{L}_n(x)$ , the set of all words of length  $n$  in  $x$ .
- Language of  $x$  is defined by  $\mathcal{L}(x) := \bigcup_{n \in \mathbb{N}} \mathcal{L}_n(x)$ .
- A substitution  $\sigma$  on  $\mathcal{A}$  is a map from  $\mathcal{A}$  to  $\mathcal{A}^+$  (the set of finite non-empty words on  $\mathcal{A}$ ), which associates to the letter  $a \in \mathcal{A}$  the word  $\sigma(a) \in \mathcal{A}^+$ , with length  $h_a := |\sigma(a)| < \infty$ .
- We define language of a substitution  $\sigma$  by :  
 $\mathcal{L}_\sigma = \{\text{factors of } \sigma^n(a) : \text{for some } n \geq 0, a \in \mathcal{A}\}$ .
- Define  $X_\sigma = \{x \in \mathcal{A}^{\mathbb{Z}} : \mathcal{L}(x) \subset \mathcal{L}_\sigma\} \subset \mathcal{A}^{\mathbb{Z}}$ .

## Preliminaries cont.

- The *left shift*  $T : \mathcal{A}^{\mathbb{Z}} \rightarrow \mathcal{A}^{\mathbb{Z}}$ , is defined by  $(Tx)_k = x_{k+1}$ , for all  $k \in \mathbb{Z}$ .
- $X_\sigma$  is a **Polish space** and is closed under  $T$ . We call  $(X_\sigma, T)$  the subshift associate with substitution  $\sigma$ .
- $(X_\sigma, T)$  is **Borel dynamical system** i.e.  $T$  is a homeomorphism of Polish space.
- For a finite string  $\bar{x} = (x_0, \dots, x_n)$  of length  $n$ , denote by  $[\bar{x}]$  a *cylinder set* :  $[\bar{x}] := \{y = (y_i) \in X_\sigma : y_0 = x_0, \dots, y_n = x_n\}$ .
- We say that  $\sigma$  is of **bounded length** if there exists an integer  $C \geq 2$  such that for every  $a \in \mathcal{A}$ ,  $|\sigma(a)| \leq C$ .  $n=1 \quad \sigma(1)$
- By identifying  $\mathcal{A}$  with  $\mathbb{Z}$ , we say that  $\sigma$  is of **bounded size**, if it is of **bounded length** and there exists a positive integer  $t$  (independent of  $n$  and minimal possible), such that for every  $n \in \mathbb{Z}$ , if  $m \in \sigma(n)$ , then  $m \in \{n-t, \dots, n, \dots, n+t\}$ .  $2t+1$

# Left determined substitution on infinite alphabet

*Definition* ([Ferenczi 2006]). We say that a substitution  $\sigma$  on a countable alphabet  $\mathcal{A}$  is **left determined** if there exists  $N \in \mathbb{N}$  such that, any word  $w \in \mathcal{L}_\sigma$  of length at least  $N$ , has a unique decomposition  $w = w_1 \dots w_s$ , such that each  $w_i = \sigma(a_i)$  for unique  $a_i \in \mathcal{A}$ , except that  $w_1$  may be a suffix of  $\sigma(a_1)$  and  $w_s$  may be a prefix of  $\sigma(a_s)$ .

$|w| = N$

$$w = \underbrace{b_1 b_2 b_3}_{\substack{\text{suffix of} \\ \sigma(a_1)}} \underbrace{b_4 b_5 b_6 b_7}_{= \sigma(a_2)} \dots \underbrace{b_{N-4} b_{N-3} b_{N-2}}_{= \sigma(a_{s-1})} \underbrace{b_{N-1} b_N}_{\text{prefix of } \sigma(a_s)}$$

Handwritten annotations:  $w_1$  and  $w_2$  have blue checkmarks above them.  $w_{s-1}$  and  $w_s$  have blue checkmarks above them. The underlines for  $w_1$  and  $w_s$  are double lines. The word "can" is written in blue below the first and last terms.

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$$w = \overbrace{b_1 b_2 b_3}^{w_1} \overbrace{b_4 b_5 b_6 b_7}^{w_2} \dots \overbrace{b_{N-4} b_{N-3} b_{N-2}}^{w_{s-1}} \overbrace{b_{N-1} b_N}^{w_s}$$

$\text{suffix of } \sigma(a_1)$        $= \sigma(a_2)$        $= \sigma(a_{s-1})$        $\text{prefix of } \sigma(a_s)$

Example: *The squared drunken man substitution:*

$$n \mapsto (n-2) \underline{\underline{nn(n+2)}}; n \in 2\mathbb{Z}$$

is left determined (see [F06]).

# Example: a (non-simple, finite rank) Bratteli diagram

$$\alpha, \beta : E \rightarrow V$$

$$\forall v \in V \setminus V_0 \quad \alpha^{-1}(v) \neq \emptyset$$

$$\forall v \in V \quad \beta^{-1}(v) \neq \emptyset$$

$$B(V, E)$$

$$|V_i| < \infty$$

$$|E_i| < \infty$$

$$V = \bigcup_i V_i$$

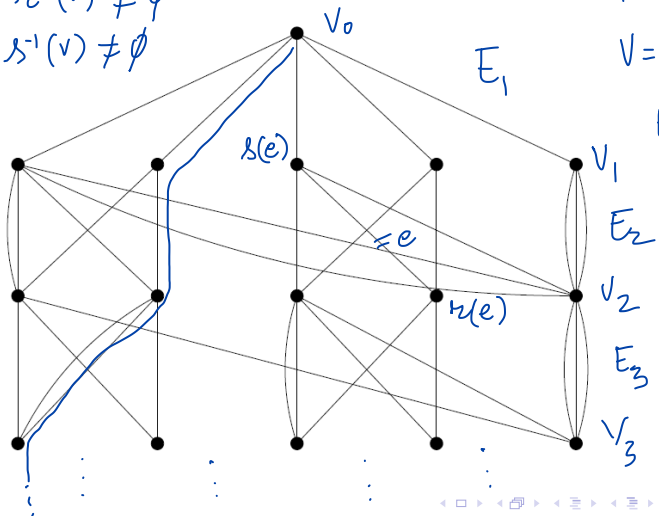
$$E = \bigcup_i E_i$$

$$(e_i)_i$$

$$\alpha(e_i)$$

$$= \beta(e_{i+1})$$

$$\forall i$$



↓  
Path space

# Definition of a Bratteli diagram

## Definition

A **Bratteli diagram** is a graded infinite graph  $B = (V, E)$  with the vertex set  $V = \bigsqcup_{i \geq 0} V_i$  and edge set  $E = \bigsqcup_{i \geq 1} E_i$ :

1)  $V_0 = \{v_0\}$  is a single point;

2)  $V_i$  and  $E_i$  are finite sets for every  $i$ ;

3) edges  $E_i$  connect  $V_{i-1}$  to  $V_i$ : there exist maps  $r$  (range) and  $s$  (source) from  $E$  to  $V$  such that  $r(E_i) \subseteq V_i$ ,  $s(E_i) \subseteq V_{i-1}$ , and  $s^{-1}(v) \neq \emptyset$ ;  $r^{-1}(v') \neq \emptyset$  for all  $v \in V$  and  $v' \in V \setminus V_0$ .

- $B$  is **stationary** if it repeats itself below the first level.
- $B$  is of **finite rank** if for all  $n \geq 1$ ,  $|V_n| \leq k$  for some positive integer  $k$ .
- We say a finite rank diagram  $B$  has **rank  $d$**  if  $d$  is the smallest integer such that  $|V_n| = d$  infinitely often.



## Definition of a Bratteli diagram (cont.)

The incidence matrix  $F_n$  is a  $|V_n| \times |V_{n-1}|$  matrix with entries

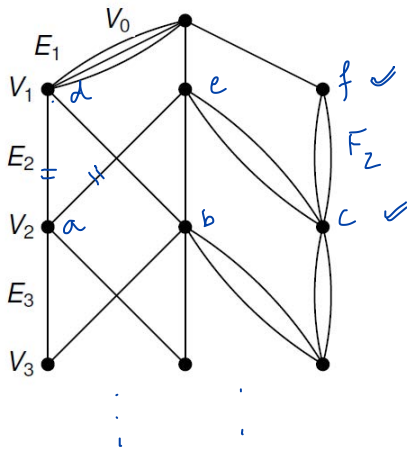
$$f_{v,w}^{(n)} = |\{e \in E_n : s(e) = w, r(e) = v\}|, \quad v \in V_n, w \in V_{n-1}. \quad \checkmark$$

A Bratteli diagram is called **simple** if  $\forall n \exists m > n$  such that  $F_m \cdots F_{n+1} > 0$  (all entries are positive).

A finite or infinite sequence of edges  $(e_i : e_i \in E_i)$  such that  $r(e_i) = s(e_{i+1})$  is called a finite or infinite path. Let  $X_B$  be the set of infinite paths starting at the top vertex  $v_0$ . Then  $X_B$  a 0-dimensional compact metric space w.r.t. the topology generated by cylinder sets

$$[\bar{e}] := \{x \in X_B : x_i = e_i, i = 0, \dots, n\}.$$

# Incidence matrix (Example)



The diagram is *stationary* with incidence matrix

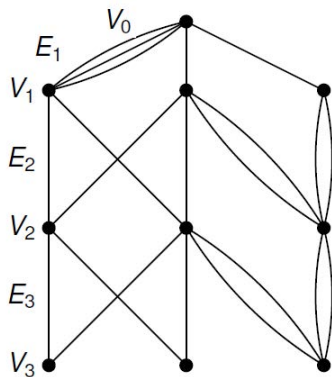
$$F = \begin{matrix} & \begin{matrix} d & e & f \end{matrix} \\ \begin{matrix} a \\ b \\ c \end{matrix} & \begin{pmatrix} 1 & 1 & 0 \\ 1 & 1 & 0 \\ 0 & 2 & 2 \end{pmatrix} \end{matrix}$$

The sequence  $(F_n)$  of incidence matrices determine the structure of a Bratteli diagram.

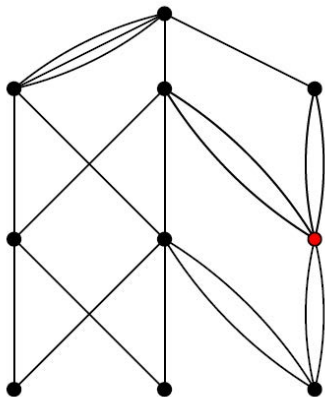
## Topology on the path space

$X_B$ : two paths are close if they agree on a large initial segment.

# Ordered Bratteli diagrams

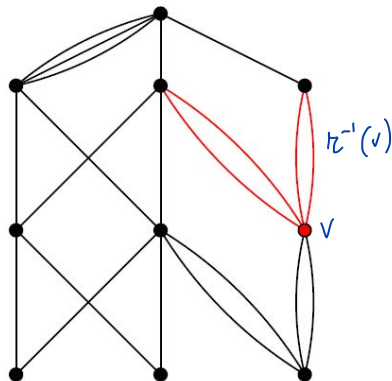


# Ordered Bratteli diagrams



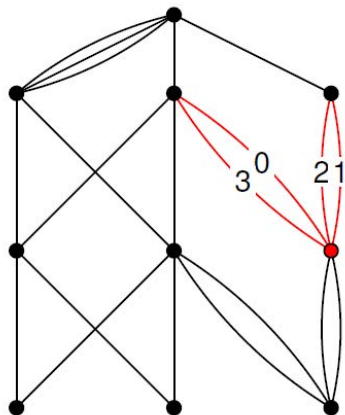
- Take a vertex  $v \in V \setminus V_0$ .

# Ordered Bratteli diagrams



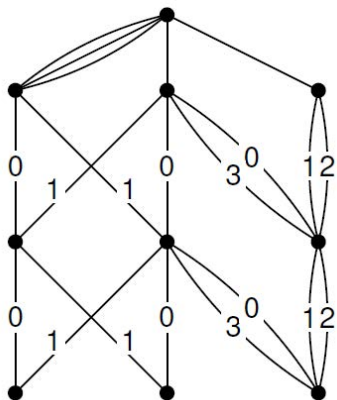
- Take a vertex  $v \in V \setminus V_0$ .
- Consider the set  $r^{-1}(v)$ .

# Ordered Bratteli diagrams



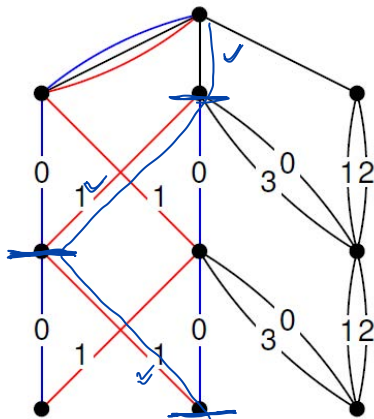
- Take a vertex  $v \in V \setminus V_0$ .
- Consider the set  $r^{-1}(v)$ .
- Enumerate edges from  $r^{-1}(v)$

# Ordered Bratteli diagrams



- Take a vertex  $v \in V \setminus V_0$ .
- Consider the set  $r^{-1}(v)$ .
- Enumerate edges from  $r^{-1}(v)$
- Do the same for every vertex.

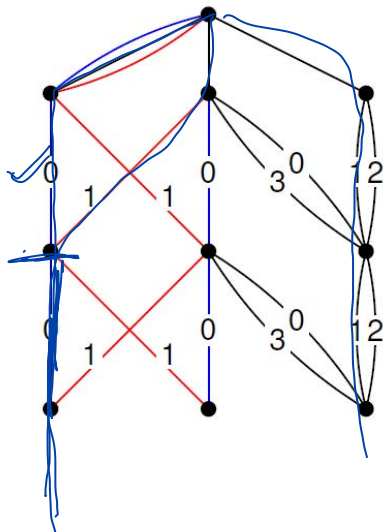
# Ordered Bratteli diagrams



- An infinite path  $x = (x_n)$  is called **maximal** if  $x_n$  is *maximal* in  $r^{-1}(r(x_n))$ . Similarly, **minimal** paths are defined.



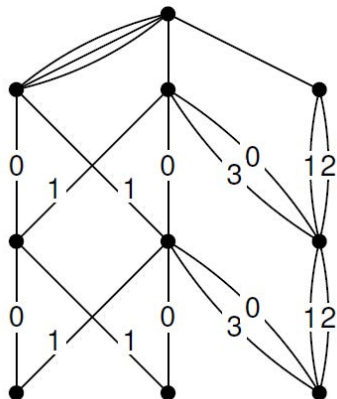
# Ordered Bratteli diagrams



- An infinite path  $x = (x_n)$  is called **maximal** if  $x_n$  is *maximal* in  $r^{-1}(r(x_n))$ . Similarly, **minimal** paths are defined.
- The sets  $X_{\max}$  and  $X_{\min}$  of all maximal and minimal paths are non-empty and closed.

# Ordered Bratteli diagrams

## Vershik map

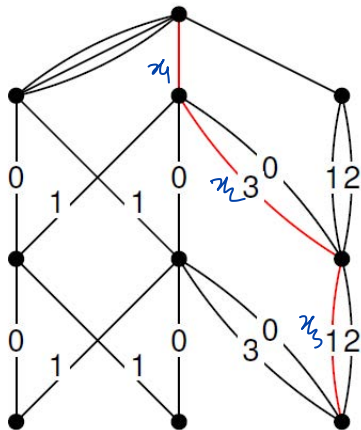


Define the **Vershik map**

$$\varphi_B : X_B \setminus X_{\max} \rightarrow X_B \setminus X_{\min} :$$

# Ordered Bratteli diagrams

## Vershik map



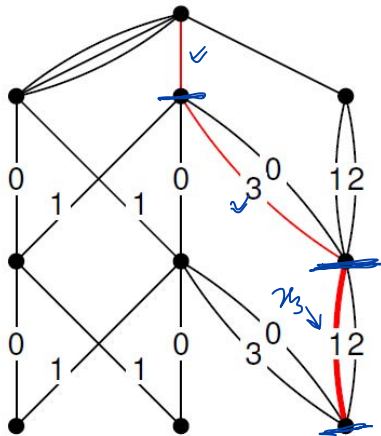
Define the **Vershik map**

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Fix  $x \in X_B \setminus X_{\max}$ .

# Ordered Bratteli diagrams

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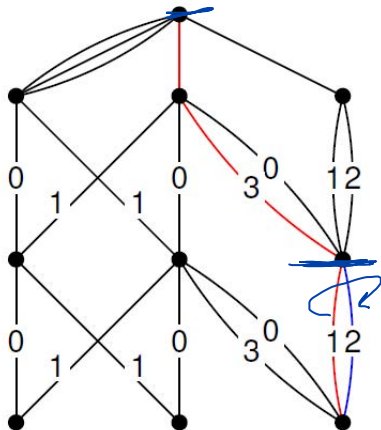
Fix  $x \in X_B \setminus X_{\max}$ .

Find the first  $k$  with non-maximal  $x_k$ .

||  
3

# Ordered Bratteli diagrams

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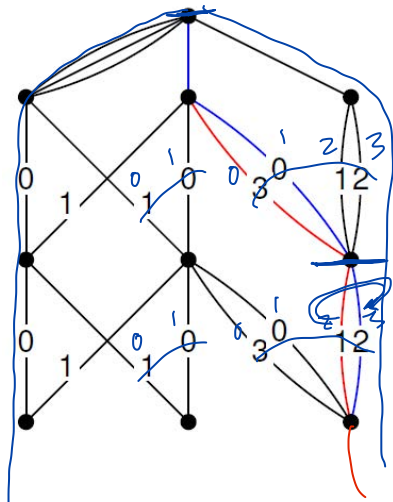
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Take  $x_k$  to its successor  $\bar{x}_k$ .

# Ordered Bratteli diagrams

## Vershik map



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Take  $x_k$  to its successor  $\bar{x}_k$ .

Connect  $s(\bar{x}_k)$  to the top vertex  $V_0$  by the minimal path.

$$\varphi_B(\text{red}) = \text{blue}$$

# Ordered Bratteli diagrams

## Vershik map

- $\varphi_B$  is defined everywhere on  $X_B \setminus X_{\max}$
- $\varphi_B(X_B \setminus X_{\max}) = X_B \setminus X_{\min}$

$\varphi_B$

### Definition

If the map  $\varphi_B$  can be extended to a homeomorphism of  $X_B$  such that  $\varphi_B(X_{\max}) = X_{\min}$ , then  $(X_B, \varphi_B)$  is called a **Bratteli-Vershik system** and  $\varphi_B$  is called the **Vershik map**.

### Question:

Under what conditions on a Bratteli diagram does the Vershik map exist?

# Ordered Bratteli diagrams

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### Answer:

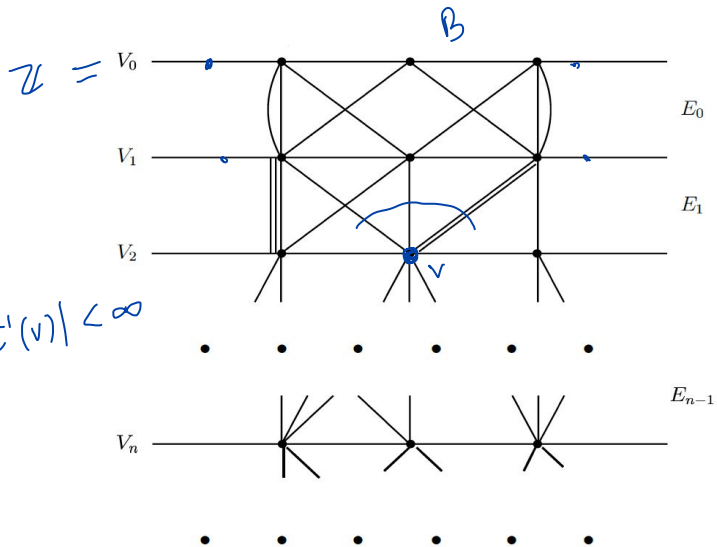
If a Bratteli diagram is **simple**, then the Vershik map **always** exists (e.g., use the left-to-right order).



## Collection of results (partial list)

- Herman, Putnam, and Skau [HPS'92] showed that for every minimal Cantor dynamical system  $(X, T)$ , there exists a simple, ordered Bratteli diagram  $B$  such that the corresponding Vershik map  $\varphi_B$  is conjugate to  $T$ .
- Bezuglyi, Dooley and Medynets (2005), Medynets (2006) extended above result to **aperiodic** Cantor dynamical systems.
- Giordano, Putnam, and Skau (1995) classified all minimal homeomorphisms of Cantor set with respect to **orbit equivalence**. By [HPS'92] it suffices to classify Vershik maps.
- Forrest (1997), Durand, Host, Skau (1999) described completely the class of dynamical systems that are represented by simple stationary Bratteli diagram. These are **minimal substitution dynamical systems**.

# Generalized Bratteli diagrams example



$|V_i|$   
 = cfb  
 infin.  
 $|E_i|$

$|\pi^{-1}(v)| < \infty$

$Y_B$

# Generalized Bratteli diagrams (GBD)

**Generalized Bratteli diagram**  $B = (V, E)$  is a countable graded graph  $B = (V, E)$  with  $V = \bigsqcup_{i \geq 0} V_i$  and  $E = \bigsqcup_{i \geq 0} E_i$  such that,

(i) The set  $V_i$ , for  $i \geq 0$  is countably infinite (identified with  $\mathbb{Z}$ ).  $E_i$  is the set of edges between the levels  $V_i$  and  $V_{i+1}$ ;

(ii) range map  $r$  and source map  $s$  from  $E$  to  $V$  such that  $r(E_i) \subset V_i$ ,  $s(E_i) \subset V_{i-1}$ ,  $s^{-1}(v) \neq \emptyset$  for all  $v \in V$ , and  $r^{-1}(v) \neq \emptyset$  for all  $v \in V \setminus V_0$ ;

(iii) for every  $v \in V \setminus V_0$ , the set  $r^{-1}(v)$  is finite. For every  $w \in V_i$ ,  $v \in V_{i+1}$ , the set of edges (denoted  $E(w, v)$ ) between  $w$  and  $v$  is finite (or empty);

(iv) Put  $|E(w, v)| = f_{vw}^i$ . This defines a sequence of infinite incidence matrices  $(F_n; n \in \mathbb{N}_0)$  whose entries are non-negative integers:

$$F_i = (f_{vw}^{(i)} : v \in V_{i+1}, w \in V_i), \quad f_{vw}^{(i)} \in \mathbb{N}_0.$$

## Generalized Bratteli diagram (GBD) cont.

- If  $F_n = F$  for all  $n \in \mathbb{N}_0$ , then diagram is called *stationary*. ✓
- We denote by  $Y_B$  the set of infinite paths in  $B = (V, E)$ .
- $Y_B$  is a **Polish space** using topology generated by cylinder sets  $[\bar{e}] := \{x \in Y_B : x_i = e_i, i = 0, \dots, n\}$ .
- $B = (V, E)$  together with a linear order on  $r^{-1}(v)$  for every  $v \in V \setminus V_0$ , is called an ordered generalized Bratteli diagram denoted by  $B = (V, E, \geq)$ .
- For an ordered diagram  $B = (V, E, \geq)$ , we define **Vershik map**  $\varphi : Y_B \rightarrow Y_B$  as before.
- $(Y_B, \varphi)$  is a **Borel** dynamical system.

### Theorem (Bezuglyi, Dooley, Kwiatkowski (2006))

Let  $T$  be an aperiodic Borel automorphism of  $(X, \mathcal{B})$ . Then there exists an ordered generalized Bratteli diagram  $B = (V, E, \geq)$  and a Vershik map  $\varphi_B : Y_B \rightarrow Y_B$  such that  $(X, T)$  is isomorphic to  $(Y_B, \varphi_B)$ .

# Back to infinite substitution : Kakutani-Rokhlin towers

For  $a_i \in \mathcal{A}$ ;  $i = \{1, \dots, n\}$ , we denote by  $[a_1 \dots a_n]$  a cylinder set of length  $n$ .

**Theorem 1.** Let  $\sigma$  be a *left determined substitution* on countably infinite alphabet  $\mathcal{A}$  and  $(X_\sigma, T)$  be the corresponding subshift. Then for any  $n \in \mathbb{N}$ , we have a partition of  $X_\sigma$  into K-R towers


$$X_\sigma = \bigsqcup_{a_1 \dots a_n \in \mathcal{L}_n(\sigma)} \bigsqcup_{k=0}^{h_1 + \dots + h_n - 1} T^k[\sigma(a_1 \dots a_n)]$$

where  $h_k = |\sigma a_k|$  for  $k \in \{1, \dots, n\}$ .

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where  $h_k = |\sigma a_k|$  for  $k \in \{1, \dots, n\}$ .

**Theorem 2.** Let  $\sigma$  be *left determined* and bounded length substitution on a countable infinite alphabet  $\mathcal{A}$ , and  $(X_\sigma, T)$  be the corresponding subshift. Then for every  $n \in \mathbb{N}$ , we have a partition of  $X_\sigma$  into K-R towers



$$X_\sigma = \bigsqcup_{a_i \in \mathcal{A}} \bigsqcup_{k=0}^{h_i^n - 1} T^k[\sigma^n a_i], \text{ where } |\sigma a_i| = h_i.$$


# Construction of Bratteli-Vershik model

We used *Theorem 1* and *Theorem 2* to obtain:

**Corollary 3.** Let  $\sigma$  be a left determined substitution of bounded size on countably infinite alphabet  $\mathcal{A}$  and  $(X_\sigma, T)$  be the corresponding subshift. Then there exists two sequence  $(A_n)$  and  $(B_n)$  of Borel sets with  $A_0 = B_0 = X_\sigma$  and for  $n > 0$ ,

$$A_n = \bigsqcup_{a_i \in \mathcal{A}} [\sigma^n a_i], \quad \text{and} \quad B_n = \bigsqcup_{a_1 \dots a_n \in \mathcal{L}_n(\sigma)} [\sigma(a_1 \dots a_n)].$$

such that

- (a)  $X_\sigma = A_0 \supset A_1 \supset A_2 \supset A_3 \dots$  and  $X_\sigma = B_0 \supset B_1 \supset B_2 \supset B_3 \dots$ .
- (b) Both  $\bigcap_{n \in \mathbb{N}} A_n$  and  $\bigcap_{n \in \mathbb{N}} B_n$  are countably infinite.
- (c)  $A_n$  and  $B_n$  are complete  $T$ -sections for each  $n \in \mathbb{N}$ .
- (d) For each  $n \in \mathbb{N}$  every point in  $A_n$  and  $B_n$  is recurrent.

# Bratteli-Vershik models for substitution on infinite alphabet

Using *Corollary 1*, (sets  $B_n$ ) we obtain Bratteli-Vershik model for left determined substitution of bounded size on countably infinite alphabet.

**Theorem 4.** Let  $\sigma$  be a left determined substitution of bounded size on an infinite alphabet  $\mathcal{A}$  and  $(X_\sigma, T)$  be the corresponding subshift. Then there exists an ordered generalized-Bratteli diagram  $B = (V, E, \geq)$  and a Vershik map  $\varphi : Y_B \rightarrow Y_B$  such that  $(X_\sigma, T)$  is isomorphic to  $(Y_B, \varphi)$ .



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Using *Corollary 1*, (sets  $A_n$ ) we obtain stationary Bratteli-Vershik model.

**Theorem 5.** Let  $\sigma$  be a left determined substitution of bounded size on an infinite alphabet  $\mathcal{A}$  and  $(X_\sigma, T)$  be the corresponding subshift. Then there exists a stationary ordered generalized-Bratteli diagram  $\tilde{B} = (\tilde{V}, \tilde{E}, \geq)$  and a Vershik map  $\tilde{\varphi} : Y_{\tilde{B}} \rightarrow Y_{\tilde{B}}$  such that  $(X_\sigma, T)$  is isomorphic to  $(Y_{\tilde{B}}, \tilde{\varphi})$ .

# Perron-Frobenius theorem for countable matrix

**Theorem 6.** (Generalized PF Theorem, see [K 98]). Suppose  $F$  is a countable, non-negative, irreducible, aperiodic matrix. Suppose that  $F$  is recurrent. Then there exists a Perron-Frobenius eigenvalue  $\lambda = \lim_{n \rightarrow \infty} (f_{ij}^{(n)})^{\frac{1}{n}} > 0$  such that:

(a) there exist unique strictly positive left  $\ell$  and right  $r$  eigenvectors corresponding to  $\lambda$ ,

(b)  $\ell \cdot r$  =  $\sum_i \ell_i r_i < \infty$  if and only if  $F$  is positive recurrent.

# Tail-invariant measure on stationary GBD

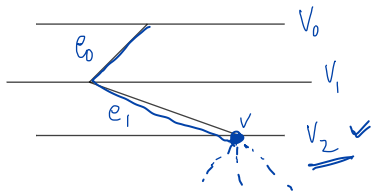
**Theorem 7.** Let  $B = B(F)$  be a stationary generalized-Bratteli diagram such that the incidence matrix  $F$ , is irreducible, aperiodic and recurrent. Then there exists a tail invariant measure  $\mu$  on the path space  $Y_B$ .

(1) Let  $[\bar{e}] = [e_0 e_1 \dots e_{n-1}]$  denote a cylinder set of length  $n$  such that  $r(\bar{e}) = v \in V_n$ , then we define :

$$\mu([\bar{e}]) = \frac{l_v}{\lambda^n}.$$

where  $l$  is the left eigenvector corresponding to Perron value  $\lambda$  of  $F$ .

Example



$$\mu([\underline{e_0 e_1}]) = \frac{l_v}{\lambda^2}$$

# Tail-invariant measure on stationary GBD

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$$\mu([\bar{e}]) = \frac{\ell_v}{\lambda^n}.$$

where  $\ell$  is the left eigenvector corresponding to Perron value  $\lambda$  of  $F$ .

(2) The measure  $\mu$  is **finite** if and only if the left eigenvector  $\ell = (\ell_v)$  has the property  $\sum_v \ell_v < \infty$ .

# From tail-invariant to shift-invariant measure

Since dynamical systems  $(X_\sigma, T)$  and  $(Y_B, \varphi)$  are isomorphic, we can push the tail-invariant measure to shift-invariant measure.

**Corollary 8.** Let  $\sigma$  be a left determined substitution of bounded size on an infinite alphabet  $\mathcal{A}$  and  $(X_\sigma, T)$  be the corresponding subshift. Assume that the countably infinite substitution matrix  $M$  is irreducible, aperiodic, and recurrent.

(1) Then there exists a shift-invariant measure  $\nu$  on  $X_\sigma$ .

(2) Let  $\ell$  be the left eigenvector of  $M$  corresponding to the Perron value  $\lambda$  of  $M$ . The measure  $\nu$  is finite if and only if the left eigenvector  $\ell = (\ell_i)$  has the property  $\sum_i \ell_i < \infty$ .

# Ergodic shift-invariant probability measure

**Theorem 9.** Let  $\sigma$  be a left determined substitution of bounded size on an infinite alphabet  $\mathcal{A}$  such that the substitution matrix  $M$  is irreducible, aperiodic, and recurrent. Then the shift-invariant **probability** measure  $\nu$  (defined in Corollary 7) on  $X_\sigma$  is ergodic.

**Proof Sketch.** We work with the stationary Bratteli-Vershik model  $(B, \varphi, \geq)$  of  $(X_\sigma, T)$ . Since  $\mu$  is probability, we have  $\sum_v \ell_v = 1$ .

$M$  is **irreducible** : For all  $i, j$  there exists some  $n$  such that  $m_{ij}^{(n)} > 0$ .

Moreover, for a fixed state  $i$  there exists  $k$  such that  $m_{ii}^{(n)} > 0$  for all  $n \geq k$ .

# Ergodic shift-invariant probability measure

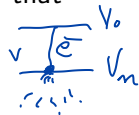
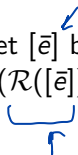
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Let  $[\bar{e}]$  be a cylinder set such that  $r(\bar{e}) = \bar{v} \in V_n$ . We will show that  $\mu(\mathcal{R}([\bar{e}])) = 1$ , where  $\mathcal{R}$  denotes the tail-equivalence relation.

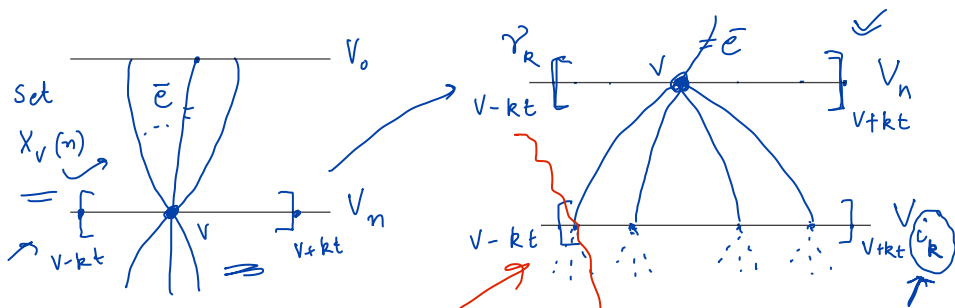


# Proof sketch cont.

For  $n > 0$  and  $v \in V_n$  we define,  $X_v(n) = \{x \in Y_B : r(x_{n-1}) = v \in V_n\}$ .

For a fixed  $k \in \mathbb{N}$  and  $t$  as in the definition of **bounded size** the set  $\mathcal{V}_k = \{v - kt, \dots, v + kt\}$  denotes a set of  $2kt + 1$  vertices.

Since  $M$  is **irreducible**, there exists a level (say  $V_{i_k}$ ,  $i_k > n$ ) such that  $v$  is connected to every  $u \in \mathcal{V}$  on level  $V_{i_k}$ .





# Proof sketch cont.

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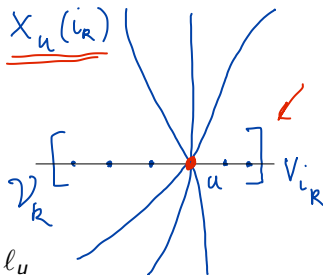
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Consider the set  $\overline{U}_k = \bigcup_{u \in \mathcal{V}_k} X_u(i_k) \subset \mathcal{R}([\bar{e}])$ .

Note that  $\mu(\overline{U}_k) = \sum_{u \in \mathcal{V}_k} \frac{|E(V_0, u)|}{\lambda^{i_k}} \ell_u$ .

$$\mu(\overline{U}_k) = \sum_{u \in \mathcal{V}_k} \frac{\sum_{w=u-i_k t}^{u+i_k t} f_{uw}^{(i_k)}}{\lambda^{i_k}} \ell_u$$



# Proof sketch cont.

By switching the order of summations

$$\mu(\overline{U}_k) = \sum_{w=u-i_k t}^{u+i_k t} \frac{\sum_{u \in \mathcal{V}_k} (\ell_u) \cdot (f_{uw}^{(i_k)})}{\lambda^{i_k}}$$

Handwritten note:  $\ell \cdot F^{i_k} = \ell \cdot \lambda$  with a red checkmark.

For  $\epsilon > 0$ , there exists  $k \in \mathbb{N}$  such that

$$\sum_{u \in \mathcal{V}_k} (\ell_u) \cdot (f_{uw}^{(i_k)}) > \underline{\lambda^{i_k} \ell_w} - \epsilon$$

$$\text{Hence, } \mu(\overline{U}_k) > \sum_{w=u-i_k t}^{u+i_k t} \frac{\lambda^{i_k} \ell_w - \epsilon}{\lambda^{i_k}} = \sum_{w=u-i_k t}^{u+i_k t} \ell_w - \sum_{w=u-i_k t}^{u+i_k t} \frac{\epsilon}{\lambda^{i_k}}$$

## Proof sketch cont.

$$\mu(\overline{U}_k) > \sum_{w=u-i_k t}^{u+i_k t} \ell_w - \frac{\epsilon}{\lambda^{i_k}} (2i_k t + 1),$$

$k \rightarrow \infty$ ,  $i_k \rightarrow \infty$ , observe that

$$\sum_{w=u-i_k t}^{u+i_k t} \ell_w \rightarrow \sum_{w \in \mathbb{Z}} \ell_w = 1 \quad \text{and} \quad \frac{(2i_k t + 1)}{\lambda^{i_k}} \rightarrow 0,$$

Thus  $\mu(\overline{U}_k) \rightarrow 1$ . Since  $\mu(\overline{U}_k) \subset \mathcal{R}([\bar{e}])$ , we get  $\mu(\mathcal{R}([\bar{e}])) = 1$ .

Thus for any two open sets  $\mathcal{O}_1, \mathcal{O}_2 \in Y_B$ , there exists  $n \in \mathbb{Z}$  such that  $\mu(\varphi_B^n(\mathcal{O}_1) \cap \mathcal{O}_2) > 0 \implies$

For any two Borel sets  $A_1, A_2$ , there exists  $n \in \mathbb{Z}$  such that  $\mu(\varphi_B^n(A_1) \cap A_2) > 0 \implies \mu$  is ergodic. □

# Summary of ongoing work

Let  $\sigma$  be a left determined substitution of bounded size on an infinite alphabet  $\mathcal{A}$  such that the substitution matrix  $M$ , is irreducible and recurrent, then following could be said about ~~the then~~ the shift-invariant measure  $\nu$  defined earlier

Substitution matrix	$\sum \nu l_\nu$	Shift-invariant measure		
		Type	Ergodic	Unique
✓ Positive recurrent	✓ Finite	✓ Finite	✓ Yes	Yes BET
Null recurrent	✓ Finite	✓ Finite	✓ Yes	??
Positive recurrent	Infinite	<u><math>\sigma</math>-finite</u>	?? ✓	If it is ergodic then it is unique. <i>Hof</i>
Null recurrent	Infinite	$\sigma$ -finite	??	??



# Example

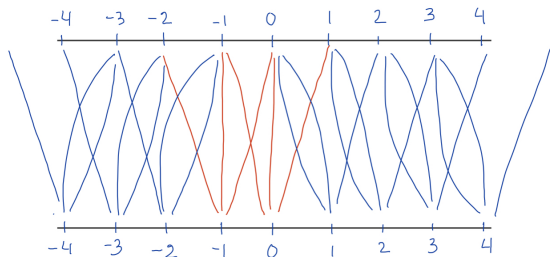
One step fwd, two step back substitution on  $\mathbb{Z}$ . Define  $\sigma$  by

$$-1 \mapsto -2 -1 0 ; \mathbf{0} \mapsto -1 0 1$$

$$n \mapsto (n-1)(n+1)(n+1) ; \underline{n \leq -2}$$

$$n \mapsto (n-1)(n-1)(n+1) ; \underline{n \geq 1}$$

This substitution is left determined, with irreducible, positive recurrent matrix;  $\lambda = 3$ ; left e.v.,  $\ell = (\dots \frac{1}{2^4}, \frac{1}{2^3}, \frac{1}{2^2}, \frac{1}{2}, 1, 1, \frac{1}{2}, \frac{1}{2^2}, \frac{1}{2^3} \dots)$



[F06] Sébastien Ferenczi. Substitution dynamical systems on infinite alphabets. Ann. Inst. Fourier (Grenoble), 56:2315-2343, 2006.

[K 98] Bruce P. Kitchens. Symbolic Dynamics. One-sided, Two-sided and Countable State Markov Shifts. Springer Universitext. ISBN 3-540-62738-3. 1998.



# Thank You

## Some definitions from theory of Stochastic Process

- Countably infinite non-negative matrix  $F$  is called *irreducible* if for all  $i, j \in \mathbb{Z}$  there exists some  $n \in \mathbb{N}$  such that  $f_{ij}^{(n)} > 0$ .
- An irreducible matrix  $F$  has period  $p$  if, for all vertices  $i \in \mathbb{Z}$ ,  $p = \gcd\{\ell : f_{ii}^{(\ell)} > 0\}$ . If  $p = 1$ , the matrix  $F$  is called *aperiodic*.
- An irreducible aperiodic matrix  $F$  admits a *Perron-Frobenius eigenvalue*  $\lambda$ , defined by,  $\lambda = \lim_{n \rightarrow \infty} (f_{ii}^{(n)})^{\frac{1}{n}}$  (independent of  $i$ ).
- An irreducible aperiodic matrix  $F$  is called *transient* if  $\sum_n f_{ij}^{(n)} \lambda^{-n} < \infty$ ; otherwise,  $F$  is called *recurrent*.
- For a recurrent matrix  $F$ , define  $\ell_{ij}(1) = f_{ij}$  and  $\ell_{ij}(n+1) = \sum_{k \neq i} \ell_{ik}(n) f_{kj}$ .
- The matrix  $F$  is called *null-recurrent* if  $\sum_n n \ell_{ii}(n) \lambda^{-n} < \infty$ ; otherwise  $F$  is called *positive recurrent*.